

Lecture 3

Instructor OH: Tuesdays 3:30-4:30 pm.

Some relevant category theory - framework

§6.2, 6.3 - but see class calendar for the topic names

(don't care what edition you have - HW will be typed out for you)

HW02 will be out tonight, or at the latest tomorrow morning

Words that I won't define in class but you're either covered or still make sense when we talk about left-modules over a non-comm ring R : (Analogous to VS)

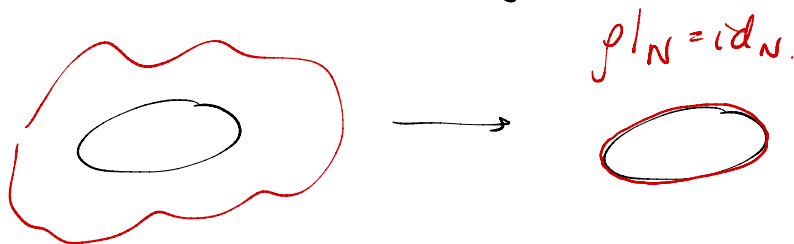
- submodule, proper submodule
- R -linear combinations
- cyclic submodule generated by $x \in M$

$$\langle x \rangle = Rx.$$

\rightsquigarrow submodule generated by a subset $X \subset M$.

- kernel, image, quotient module, cosets
- injection, projection, direct sum
 (?) cat theoretically, \times, \oplus ?
- retraction? $N \subseteq M$ as modules.

$p: M \rightarrow N$ is a retraction if $p(n) = n \quad \forall n \in N$.



N is a retract of M .

- exact sequence of modules

$$0 \longrightarrow A \xrightarrow{f} B \xrightarrow{g} C \longrightarrow D$$

(extension of A by C)

will cover next time if not clear.

$R = \text{ring}$. $S, T \in {}_R \text{Mod}$.

We wish to define \oplus, \times
in such a way that we can
generalize beyond a finite #
of factors.

+

Let \mathcal{C} be a category.

Recall Defn A morphism $f: A \rightarrow B$ in \mathcal{C} is an
isomorphism if there exists a morphism $g: B \rightarrow A$ in \mathcal{C}

where

$$\underbrace{gf = 1_A}_{\text{Mor}_{\mathcal{C}}^{\circ}(A, A)} \quad \text{and} \quad \underbrace{fg = 1_B}_{\text{Mor}_{\mathcal{C}}^{\circ}(B, B)}$$

Then g is called the inverse of f .

as opposed to left- or right inverse, which are
weaker!

"clear": • identity ^{aka homs} morphisms are always isomorphisms.
• Inverses are unique (why?)

Q. What categories are you already comfy with?

Set, ComRing, Groups, ${}_R \text{Mod}$, $\text{Mod } R$, Top?

(for examples later)

Coproducts: Cat'l.

defn. Let $A, B \in \text{Ob}(\mathcal{C})$. Their coproduct $A \sqcup B$ ($A \sqcup B$)

(ie a notion of disjoint union) is an object $C \in \text{Ob}(\mathcal{C})$

together w/ injections $\alpha: A \rightarrow A \sqcup B$

$$\beta: B \rightarrow A \sqcup B$$

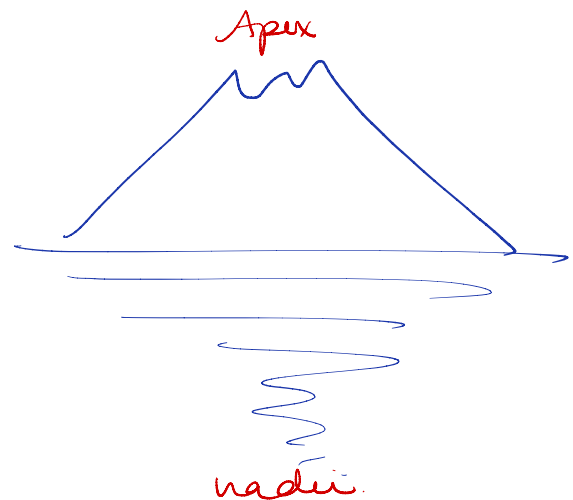
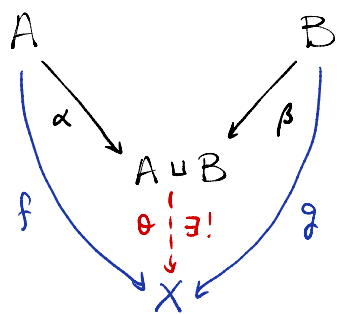
such that, $\forall X \in \text{Ob}(\mathcal{C})$

and morphisms $f: A \rightarrow X, g: B \rightarrow X$

will write $X \in \mathcal{C}$
henceforth

$\exists!$ morphism $\theta: A \sqcup B \rightarrow X$

making the diagram below commute:



eg. Coproduct in Sets, Top?

ex. What is coproduct in Groups?

prop. If $A, B \in R\text{Mod}$, then $A \perp B$ exists in $R\text{Mod}$ a coproduct (the coproduct)

and this is the direct sum $C = A \oplus B$ we expect. key word of this proof.

What do we expect? From prev class or analogy...

$$C = A \times B \text{ as set, } \begin{array}{l} A \hookrightarrow C \\ a \mapsto (a, 0) \end{array} \quad \begin{array}{l} B \hookrightarrow C \\ b \mapsto (0, b) \end{array}$$

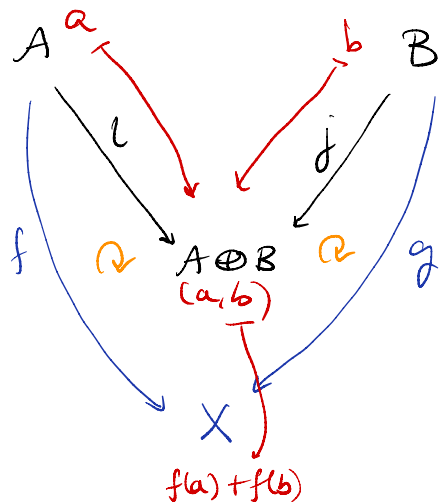
addition by coordinate: $(a, b) + (a', b') = (a+a', b+b')$

scalar mult by coord: $r(a, b) = (ra, rb)$

(can check distributivity easily)

Pf. θ exists:

θ is unique: If ψ were some other module hom,



$$\psi((a, 0)) = f(a)$$

$$\psi((0, b)) = g(b)$$

ψ is a hom
 $\Rightarrow \psi((a, b)) = \psi((a, 0)) + \psi((0, b))$

$$f(a) + g(b) = \theta(a, b)$$

□

Similar proof shows coproducts exist in Mod_R .

"Universal Property" A categorical construction (eg. coproduct) X can be defined by a universal property if whenever X exists, it is unique up to unique isom.

This is not a helpful defn — use examples.

(If you want a rigorous defn, use categorical language)

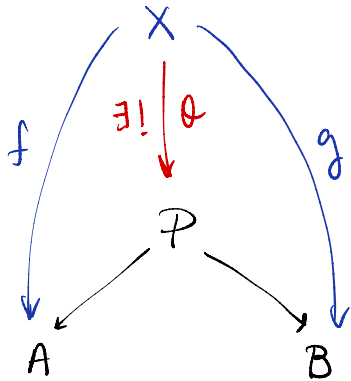
ex. Prove that if a coproduct exists, it is unique up to unique isom.

Solution is in the book... Might also be on HW in some form.

Dual notion: Products (Cat'1)

defn. Let $A, B \in \mathcal{C}$. Then their product $A \amalg B$ is an object $P \in \text{Ob}(\mathcal{C})$ along with projections $p: P \rightarrow A, q: P \rightarrow B$ such that
(blah blah, let's draw a picture)

$\prod (x, y) \dots$
notation...
I'll just
write
 $A \times B \dots$



Commutates.

prop. In $\mathcal{R}\text{Mod}$, products exist and are unique up to unique isom.

pf. HW!

Ex. ① What are $A \amalg B$ and $A \amalg B$ in Sets?

② Show these are not ism notions.

Ex. Show $A \oplus B \cong A \times B$ for $A, B \in \mathcal{R}\text{Mod}$

(and thus also $\text{Mod } \mathcal{R}$).

pf. in both... maybe also HW.

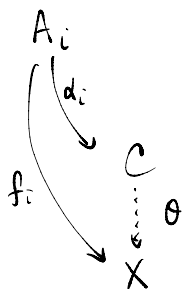
When the difference matters: infinite indexing sets

defn. $(A_i)_{i \in I}$ $I = \text{index set}$, $A_i \in \mathcal{C}$.

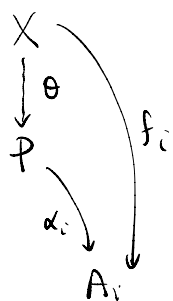
"family of objects" indexed by a set I .

Their coproduct is (the data) $(C, \{\alpha_i: A_i \rightarrow C\})$

such that $\exists! \theta$ s.t. $\forall i$, this commutes:



defn Similar for product:



Cor of prev. prop: $\bigoplus A_i$ and $\prod A_i$ are isom when
 I is finite (by induction)

claim \bigoplus and \prod are not the same in $R\text{-Mod}$! (Same for finite I by induction)

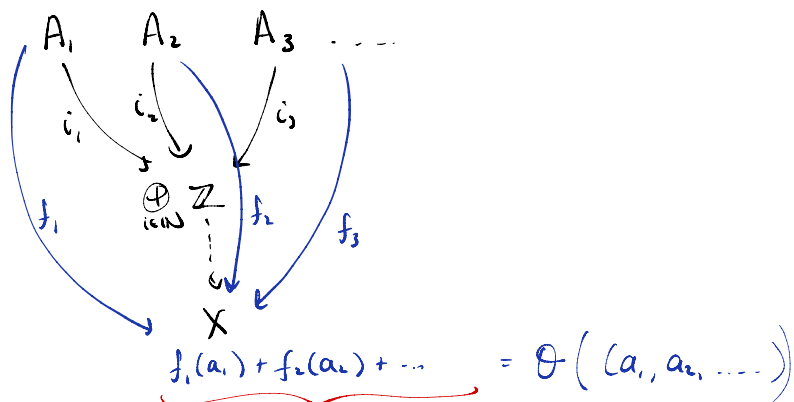
eg. Consider $\mathbb{Z}\text{-Mod} = \text{Ab}$

$$\bigoplus_i A_i = ?$$

$$\prod_i A_i = ?$$

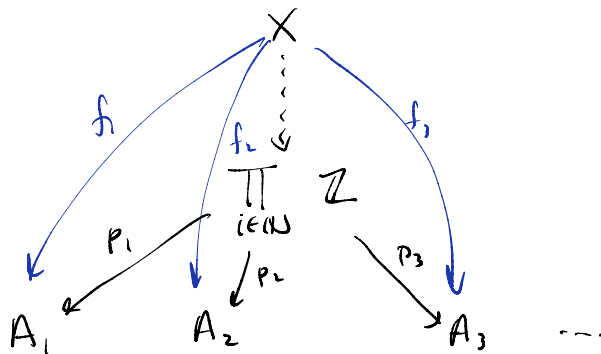
eg. let all $A_i = \mathbb{Z}$. let $I = \mathbb{N}$.

We can only add finitely many things in $\mathbb{Z} \Rightarrow$



only makes sense in $\mathbb{Z}\text{-mod } X$ if we have a finite sum.

eg. On the other hand:



no problem ... we can have ∞ many nonzero entries.