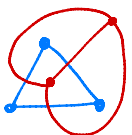


LECTURE 7

Announcements: I'm giving a 2-part lecture series @ Mathclub + AUM this week. M 5:30 - 6:30 pm in MSB Colloquium Room (MSB 1147)

defn. The planar dual of a planar graph T is given by

- $V =$ regions of the plane cut out by T
- $E =$ edge b/w v_1 and v_2 for every edge in T that is part of the border between them

eg.  the blue and red planar graphs are dual to each other

Q let D be a knot diagram. View D as a graph where

- vertices = crossings
- edges = arcs of the knot diagram

What do you notice about the planar dual of D ?

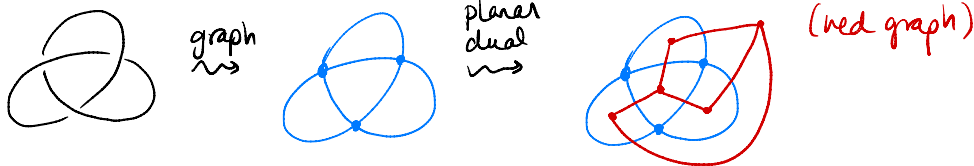
A. Since each vertex of D has valence 4, the dual graph

↑ # edges coming out of it

is composed of 4-cycles (ie squares \square),

ie the boundary of each region of the dual graph is a 4-cycle.

eg.



Sketch of proof of checkerboard coloring existence

- dual graph made of squares
- color vertices along spanning tree
- verify that adding edges doesn't cause obstruction

prop. Every knot diagram D has a checkerboard coloring.

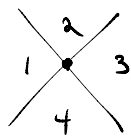
Proof sketch

Let D be a knot diagram; let Γ be the dual graph to D .

Claim A Γ is composed of 4-cycles glued together along edges.

Pf of Claim

- each vertex x (crossing) of D is 4-valent.
- the four regions touching at a crossing must all be distinct



- there is a 4-cycle in Γ surrounding each vertex
- these cycles are glued along their edges, wherever D had an edge connecting two crossings



because of the blue edge in D , the purple arcs in the left picture are actually the same edge in the dual graph Γ .

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It suffices to color the vertices of Γ 2 colors, so that if $(v_1, v_2) \in E(\Gamma)$, then v_1 and v_2 are colored differently.

Algorithm: Pick a vertex and color a random color. Then propagate. There is only a constraint if you travel along a loop (ie cycle) to a vertex you've already colored.

But by the claim, every cycle in Γ is even length

(prove this). So your coloring will be consistent. \square

Knot Polynomials via skein relations

Alexander Polynomial (1923)

$\Delta_K(t) \in \mathbb{Z}[t, t^{-1}]$ Laurent polynomials, K oriented

Defined recursively:

- $\Delta_U(t) = 1$
- $\Delta_{L_+}(t) - \Delta_{L_-}(t) = (t^{1/2} - t^{-1/2}) \Delta_{L_0}(t)$

i.e.



local pictures

$$\langle \begin{array}{c} \nearrow \\ \searrow \\ L_+ \end{array} \rangle - \langle \begin{array}{c} \nwarrow \\ \nearrow \\ L_- \end{array} \rangle = (t^{1/2} - t^{-1/2}) \langle \begin{array}{c} \uparrow \\ \uparrow \\ L_0 \end{array} \rangle$$

shorthand.

eg. Resolving tree:



Fact $\Delta_K(t^{-1}) = \Delta_K(t)$ for all knots (symmetric / palindromic)

Pf. look at skein relation.